

### Water Resources Systems: Modeling Techniques and Analysis

Lecture - 29

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#### Summary of the previous lecture

#### Normal distribution

$$f(x) = \frac{1}{\sqrt{2\pi}\sigma} \exp\left\{-\frac{1}{2} \left(\frac{x-\mu}{\sigma}\right)^2\right\} - \infty < x < +\infty \text{ Symmetric about } x = \mu$$

$$\sum_{x < x < +\infty} \int_{\mu}^{\pi} \int_{x \to \infty}^{\pi} \int_{$$

$$Z = \frac{X - \mu}{\sigma}$$

$$f(z) = \frac{1}{\sqrt{2\pi}} e^{-z^2/2}$$

$$-\infty < z < +\infty$$

$$Z = \frac{X - \mu}{\sigma} \qquad f(z) = \frac{1}{\sqrt{2\pi}} e^{-z^2/2} \qquad -\infty < z < +\infty \qquad \text{N(M, 2)}$$
• Lognormal Distribution
$$f(x) = \frac{1}{\sqrt{2\pi}x\sigma_y} e^{-(\ln x - \mu_y)^2/2\sigma_y^2} \qquad 0 < x < \infty, 0 < \mu_y < \infty, \sigma_y > 0$$

$$f(x) = \frac{1}{\sqrt{2\pi}x\sigma_y} e^{-(\ln x - \mu_y)^2/2\sigma_y^2}$$

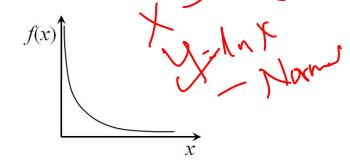
$$0 < x < \infty, 0 < \mu_y < \infty, \sigma_y > 0$$

**Exponential Distribution** 

$$f(x) = \lambda e^{-\lambda x}$$

$$x > 0, \lambda > 0$$

$$f(x) = \lambda e^{-\lambda x} \qquad x > 0, \lambda > 0$$
$$F(x) = 1 - e^{-\lambda x} \qquad x > 0, \lambda > 0$$



## CHANCE CONSTRAINED LINEAR PROGRAMMING (CCLP) FOR RESERVOIR DESIGN AND OPERATION

Deterministic LP model

Min Ks.t.  $S_{t+1} = S_t + Q_t - X_t - R_t$   $R_t \ge D_t$   $S_t \le K$   $R_t \le R_t^{\max}$ 

 $Q_t$ , inflow during period, t, is a random variable.

Probability distribution of  $Q_t$  is known

# Storage, $S_t$ Inflow, $Q_t$ Release, $R_t$ Known demand, $D_t$ Random variable Deterministic

 $S_t$  and  $R_t$  being functions of  $Q_t$ , are also random in nature

In a constraint containing two rvs, if the probability distribution
of one is known, the probabilistic behavior of the second can
be expressed as a measure of probability in terms of the
probability of the first variable.

#### Chance constraint:

The constraint relating release,  $R_{t}$  (random) and demand,  $D_t$  (deterministic) is expressed as a 

$$P[R_t \geq D_t] \geq \alpha_1$$

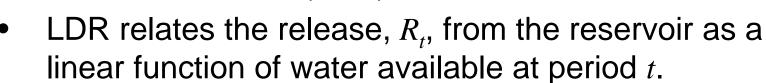
Probability of release equaling or exceeding the known demand is at least equal to  $\alpha_1$  .... referred as reliability level.

The reliability of meeting the demand in period t is at least  $\alpha_1$ .

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- Similarly,  $P[R_t \le R_t^{\max}] \ge \alpha_2$   $P[S_t \le K] \ge \alpha_3$   $P[S_t \ge S_{\min}] \ge \alpha_4$
- Probability distribution of  $S_t$  and  $R_t$  to be determined from known probability distribution of  $Q_t$ .
- Since  $S_t$ ,  $R_t$  and  $Q_t$  are interdependent, it is not possible to derive both probability distributions of  $S_t$  and  $R_t$ .
- To overcome this difficulty, Linear Decision Rule (LDR) is appropriately defined.

#### Linear Decision Rule (LDR):



$$R_{t} = S_{t} + Q_{t} - b_{t}$$

 $R_{t} = S_{t} + Q_{t} - b_{t}$   $b_{t} \text{ is a deterministic parameter (decision)}$ parameter).

- In this LDR, the entire amount,  $Q_t$ , is taken into account while making release decision.
- Depending on the proportion of inflow,  $Q_t$ , used in the LDR, a number of such LDRs may be formulated.

A general form of LDR may be written as

$$R_t = S_t + \beta_t Q_t - b_t \qquad 0 < \beta_t < 1 \qquad \forall t$$

- $\beta_t = 0$  yields a relatively conservative release policy with release decisions related only to the storage.
- $\beta_t = 1$  yields an optimistic policy where the entire amount of water is available  $(S_t + Q_t)$ , is used in the LDR.

Consider the LDR

$$R_t = S_t + Q_t - b_t$$

Storage continuity equation is

$$S_{t+1} = S_t + Q_t - R_t$$

uity equation is  $S_{t+1} = S_t + Q_t - R_t$ 

results in  $S_{t+1} = b_t$ 

$$S_{t+1} = b_t$$

 $S_{t+1}$ , is set equal to  $b_t$ 

- Treat  $S_t$  deterministic in the formulation.
- Advantage: other rv,  $R_t$ , may be expressed in terms of known distribution of  $Q_t$ . (Variance of  $Q_t$  is entirely transferred to the variance of  $R_t$ )

Deterministic constraint of a chance constraint:

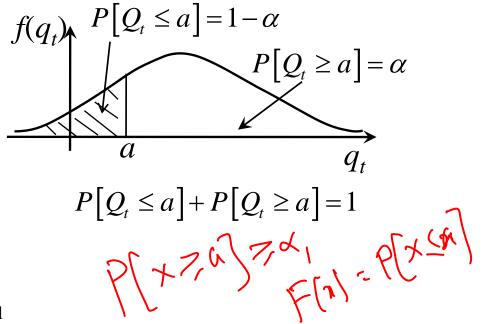
$$P[R_{t} \geq D_{t}] \geq \alpha_{1}$$

$$P[S_{t} + Q_{t} - b_{t} \geq D_{t}] \geq \alpha_{1}$$

$$P[b_{t-1} + Q_{t} - b_{t} \geq D_{t}] \geq \alpha_{1}$$

$$P[Q_{t} \geq D_{t} + b_{t} - b_{t-1}] \geq \alpha_{1}$$

$$P[Q_{t} \leq D_{t} + b_{t} - b_{t-1}] \leq 1 - \alpha_{1}$$



Deterministic with  $b_{t-1}$  and  $b_t \to \text{decision variables}$  and  $D_t \to \text{known demand}$ 

$$P[Q_{t} \leq D_{t} + b_{t} - b_{t-1}] \leq 1 - \alpha_{1}$$

Equation rewritten as

$$F_{\mathcal{Q}_t}\left(D_t + b_t - b_{t-1}\right) \le 1 - \alpha_1$$

 $F_{\mathcal{Q}_t}\left(D_{t}+b_{t}-b_{t-1}
ight)$  denotes CDF of the rv  $\mathcal{Q}_t$ 

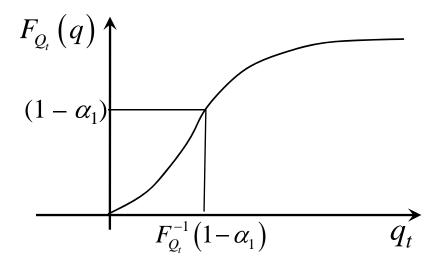
Deterministic equivalent of chance constraint  $P[R_t \ge D_t] \ge \alpha_1$  is rewritten as

$$(D_{t} + b_{t} - b_{t-1}) \leq F_{Q_{t}}^{-1} (1 - \alpha_{1})$$

- De terministic Equivalent

 $F(a) = P(x \leq a)$ 

 $F_{\mathcal{Q}_t}^{\scriptscriptstyle -1}(1-lpha_{\scriptscriptstyle 1})$  is the flow,  $q_t$  , at which the CDF value is  $1-lpha_{\scriptscriptstyle 1}$ 



The deterministic equivalent of a chance constraint,  $P \lceil R_t \le R_t^{\text{max}} \rceil \ge \alpha_2$ , is similarly obtained as

$$\left(R_t^{\max} + b_t - b_{t-1}\right) \ge F_{Q_t}^{-1}\left(\alpha_2\right)$$

The other two constraints

$$P[S_{t} \leq K] \geq \alpha_{3}$$

$$P[S_{t} \geq S_{\min}] \geq \alpha_{4}$$

 $P[S_t \le K] \ge \alpha_3$   $P[S_t \le S_{\min}] \ge \alpha_4$ Become deterministic constraints since the storage is set equal to the deterministic parameter  $b_{t-1}$ 

The complete deterministic equivalent of CCLP is written as

S.t. 
$$(D_t + b_t - b_{t-1}) \leq F_{Q_t}^{-1} (1 - \alpha_1)$$

$$(R_t^{\max} + b_t - b_{t-1}) \geq F_{Q_t}^{-1} (\alpha_2)$$

$$b_{t-1} \leq K$$

$$b_{t-1} \geq S_{\min}$$

$$b_t \geq 0$$

$$K \geq 0$$